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#### Streaming Molecular Simulation: Gromacs, OpenMM & CUDA

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# Comme test

#### namics MD on GPUs Is The good, the bad, and the ugly



Should we change the way we do HPC?



Ensemble simulations

# Molecular Dynamics on traditional CPUs

### Biomolecular Dynamics



#### **Protein Folding**

Understand biology

### Free Energy & Drug Design

#### **Timescales of Motion**



## Modeling levels

- QM
- Car-Parinello dynamics
- Molecular mechanics simulations
- Mesoscopic models
- Which is more important?
  - More detailed model
  - Better phase-space sampling

#### Detailed

Better sampling

### Molecular Mechanics



One small step for H<sub>2</sub>O





...one giant leap for water!

if (nsr[jgid] >= MAX\_CG) {
 put\_is\_list(bHaveL7,ngid,m
 nsr(jgid],n) ert





#### 20 years of CPU Optimization

- Single precision when possible
- Handtuned 1/sqrt(x) instructions
- Single-instruction Multiple-Data
- Don't calculate zeros
  - No charge on an atom? Don't do coulomb!
  - Interaction specific kernels
  - Complex neighbor lists
- Fancy algorithms to extend timestep
  - Bond constraints, virtual sites, etc.

1 nlist entry 9 interactions

"Save FLOPS

#### Stuff that might be hard to port to Cuda (I):

### Constraints

#### • Δt limited by fast motions - 1fs

Remove bond vibrations

#### • SHAKE (iterative, slow) - 2fs

- Problematic in parallel (won't work)
- Compromise: constrain h-bonds only -1.4fs

#### • GROMACS (LINCS):

- LINear Constraint Solver
- Approximate matrix inversion expansion
- Fast & stable much better than SHAKE
- Non-iterative
- Enables 2-3 fs timesteps
- Parallelizes (in theory at least)

LINCS: t=2' t=1 A) Move w/o constraint

> t= B) Project out motion along bonds

C) Correct for rotational

extension of bond

#### Stuff that might be hard to port to Cuda (2):

### Virtual sites

- Next fastest motions is H-angle and rotations of CH<sub>3</sub>/NH<sub>2</sub> groups
- Try to remove them:
  - Ideal H position from heavy atoms.
  - CH<sub>3</sub>/NH<sub>2</sub> groups are made rigid
  - Calculate forces, then project back onto heavy atoms
  - Integrate only heavy atom positions, reconstruct H's
- Enables 5fs timesteps!

a a lbl



Problem-specific optimization: use our knowledge of chemistry





#### **CPU Parallelization**



#### Complex Load balancing

CPU limit today: 50-150 ns/day

Hard to decrease # atoms / CPU



Lots of book-keeping

# Molecular Dynamics on Stream Processors

### Stream computing



Our first cards were less fancy!

First Gromacs GPU project in 2002 with Ian Buck & Pat Hanrahan, Stanford Promise of theoretical high FP performance on GeForce4 Severe limitations in practice...

#### Lessons:

It's easy to achieve speedup relative to a slow reference implementation Much harder to beat well-optimized CPU code that can use >1 core Don't write CPU code for GPUs Find new optimizations instead

#### Molecular Dynamics on GPUs



Focus on the code where we spend all the cycles on the CPU

Bottlenecks



#### OpenMM

Not sufficient to accelerate nonbonded interactions - need to send to/from GPU

Need to do entire simulation on GPU?

Not fun to rewrite 2M lines-of-code in a separate CUDA-Gromacs...



OpenMM: Core MD functionality in separate library Stanford (Pande), Stockholm (Us), Nvidia & AMD Fully public API, hardware-agnostic, use anywhere Peter Eastman, Scott Legrand

### Nonbonded Interactions



Divide into short-range / long-range

- Calculate short-range analytically
- Use approximations for the long-range

### O(N<sup>2</sup>) Algorithm on GPUs

An efficient algorithm should:
minimize global memory access
avoid thread synchronization
take advantage of symmetry

### O(N<sup>2</sup>) Algorithm on GPUs

- Group atoms into blocks of 32
- Interactions divide into 32x32 tiles
- Each tile is processed by a group of 32 threads
  - Load atom coordinates and parameters into shared memory
  - Each thread computes interactions of one atom with 32 atoms
  - Use symmetry to skip half the tiles



### O(N<sup>2</sup>) Algorithm on GPUs

- Each thread loops over atoms in a different order
  - Avoids conflicts between threads
- No explicit synchronization needed
  - -Threads in a warp are always synchronized



### O(N) Algorithm on GPUs

#### Hard. Why?

- Traditional O(N) methods (e.g. neighbor lists) slow on GPUs
  - Out of order memory access
    - for i = 1 to numNeighbors slow!
       load coordinates and parameters for neighbor[i]
       compute force slow!
       store force for neighbor[i]

Inner loop contains non-coalesced memory access

### Approach 1: Voxels

- Divide space into smaller voxels
- Compute all-vs-all interactions between adjacent voxels
- LAMMPS, NAMD
- Very efficient when all particles are identical
- Book-keeping can get expensive for complex systems



#### Approach 2: Tiles

- Start with the O(N<sup>2</sup>) algorithm
- Exclude tiles with no interactions
  - Like a neighbor list between blocks of 32 atoms

Problem 1: How do you construct good tiles?

Problem 2: How do you keep them good?



#### Finding tiles with interactions

- Compute an axis aligned bounding box for the 32 atoms in each block
- Calculate the distance between boxes



### Keeping track of water

- Solvent molecules must be ordered to be spatially coherent
  - Or bounding boxes will be very large
- Arrange along a space filling curve
- Reorder every ~100 time steps

Swapping is easy: Same parameters!



#### Performance?

- Much faster than O(N<sup>2</sup>) for large systems
- Performance scales linearly

#### But

- Computes many more interactions than really required
  - Computes all 1024 interactions in a tile, even if few/none are within the cutoff

### Finer-grained neighborlist

- For each tile with interactions:
  - Compute distance of each atom in one block from the bounding box of the other block
  - Set a flag for each atom



### **Tile Force Computation**

for i = 1 to 32

if (hasInteractions[i])

compute interaction with atom i

All 32 threads must loop over atoms in the same order

- Requires a reduction to sum the forces
- For a few atoms, this is still much faster
- -For many atoms, better to just compute all interactions



on each thread

### PME (Long range coulomb)



- 1. Calculate scaled fractionals
- 2. Calculate B-spline coefficient
- 3. "smear" the charges over the grid points from spline-coefficient
- 4. Execute forward FFT
- 5. Calculate the reciprocal energy
- 6. Execute backward FFT
- 7. Calculate force gradient

#### PME (Long range coulomb)



GPU gathering of charges





sort the atoms before gather !

#### Gromacs & OpenMM in practice

- GPUs supported in Gromacs 4.5 mdrun ... -device "OpenMM:Cuda"
- Same input files, same output files: "It just works"
- Subset of features work on GPUs for now (checked)
- No shortcuts taken on the GPU:
  - At least same accuracy as on the CPU (<1e-6)
  - Potential energies calculated, free energy works
- Prerelease availability: NOW! www.gromacs.org/gpu

#### Fermi (C20) performance over C10 BPTI (~21k atoms) Villin (600 atoms, implicit)





![](_page_36_Figure_0.jpeg)

#### Limitations

- Still hard to use long time steps on GPUs
  - Virtual sites don't work
  - We don't think you should go higher than 2fs steps without them
- Many of the CPU "tricks" can easily be written in Cuda, but we would end of with lots of kernels that must be called iteratively
- Hard to get multi-node GPU code to beat CPUs
  In the high end, we're all bandwidth-limited

#### Hardware Caveats

- Beware of Memory Errors: happens on all hardware
- Gromacs runs tests to check for GPU memory errors
  - Low-end consumer cards can sometimes be bad
- Even fine cards can exhibit random errors
- For production scientific work you might want Tesla-class Fermi cards...
  - Why? ECC memory! (C2050/C2070)

![](_page_38_Picture_7.jpeg)

#### GPU weak scaling

![](_page_39_Figure_1.jpeg)

d) >= HAX\_CG) {
 ss(bHaveLJ,ngid,m
 nsr(jgid],n1 sr)

#### More GPU/CPU comparisons

Nvidia Tesla 2050 vs Intel Core i7 920 (2.66 GHZ, 4 threads)

![](_page_40_Figure_2.jpeg)

real space cost  $\sim r_c^3$ FFT cost  $\sim$  spacing<sup>-3</sup> for constant accuracy: spacing =  $r_c/a$ total cost:  $C_{pp} r_c^3 + C_{FFT} a^3 r_c^{-3}$ 

### Tiling circles is difficult

#### serial computing

![](_page_41_Figure_2.jpeg)

#### stream computing

![](_page_41_Figure_4.jpeg)

You need lots of cubes to cover a sphere
 All interactions beyond cut-off must be zero

### The Art of Calculating Zeros

![](_page_42_Figure_1.jpeg)

# Parallelize the Problem

![](_page_44_Picture_0.jpeg)

#### These will soon be small computers ~2024: 1B'cores' 2022: ~300M cores 2020: ~100M cores 2018: ~30M cores 2016: ~10M cores 2014: ~3M cores How will YOU 2012: ~1M cores use a billion cores? 2010: ~300,000 cores

#### We're all doing Embarrassing Parallelism But not the way you think.

We're investing huge efforts in parallelizing algorithms that only reach 50-75% scaling efficiency on large problems

Not a chance they will scale to 1B cores

**Close-to-useless for smaller problems of commercial interest** 

100% focus on programs, forget the problem we're solving

Ask taxpayers to foot the bill

Pretty much the definition of 'embarrassing'?

### Scaling as an Obsession?

![](_page_46_Figure_1.jpeg)

Gromacs has scaled to 150k cores on Jaguar @ ORNL

Only gigantic systems scale - limited number of applications

![](_page_46_Picture_4.jpeg)

![](_page_46_Picture_5.jpeg)

1M-100M atoms But: <u>Small</u> systems won't scale to large numbers of cores! How can we break this impasse?

#### **Ensemble Simulation**

![](_page_47_Picture_1.jpeg)

#### Membrane Protein Insertion Free Energy

![](_page_47_Figure_3.jpeg)

#### 

### Vesicle fusion - 1.5M atoms

![](_page_48_Picture_1.jpeg)

### Long fusion trajectories

- Run on Infiniband cluster, ~250-500 cores
- 7 vesicle pairs fusing in 100-250ns
  - 75-100% POPE lipids
- 2 non-fusing vesicle pairs, 100ns & 500ns
  - 50% POPE lipids
- Interesting circumstantial observations, but hard to draw conclusions from

#### **Committer Analysis**

- Pick 20 conformations along fusion path
- Restart with 20 different random seeds
- 'Shooting trajectories'
- 8 microseconds of additional simulation
- Run on capacity cluster, 16-32 cores each
- Folding@Home as cluster scheduler
- Calculate fusion probabilities

#### Stalk commitment

![](_page_51_Figure_1.jpeg)

#### Statistics from 1000's of runs

![](_page_52_Figure_1.jpeg)

#### Scale the Problem, not Runs

![](_page_53_Picture_1.jpeg)

- Stream Computing is the future for all HPC
- We're doing *statistical* mechanics!
- No algorithm will parallelize 5000 degrees of freedom over 1 billion processors
- Parallelize in the problem domain instead
- Node efficiency becomes the key measure

#### Acknowledgments

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- Ensemble Simulations: Peter Kasson

![](_page_54_Picture_5.jpeg)