# Streaming Molecular Simulation: Gromacs, OpenMM \& CUDA 

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Molecular Dynamics on CPUs


MD on GPUs
The good, the bad, and the ugly


Should we change the way we do HIPC?


Ensemble simulations

# Molecular Dynamics 

 on traditional CPUs
## Biomolecular Dynamics



## Timescales of Motion

## Experiments



Efficient averaging
Less detail

## Where we



## Modeling levels

Detailed

- QM
- Car-Parinello dynamics
- Molecular mechanics simulations
- Mesoscopic models

Better sampling

- Which is more important?
- More detailed model
- Better phase-space sampling


## Molecular Mechanics

## Bond vibration



Torsion potentials

van der Waals interactions


Electrostatics


## One small step for $\mathrm{H}_{2} \mathrm{O}$.

8 fs


## ...one giant leap for water!

## 10 ps!

## Most interesting systems are not homogenous...

## $\partial^{2} r_{i}$ <br> $$
m_{i} \frac{\partial r_{i}}{\partial t^{2}}=F_{i} \quad i=1 . . N
$$

$$
F_{i}=-\frac{\partial V(r)}{\partial r_{i}}
$$

mimporex

$$
+\sum_{i j} \frac{q q q_{j} i_{i j}}{4 \pi e_{0}}
$$

$$
+\sum_{i, j}\left[\frac{c_{12}}{r_{12}}-\frac{C_{6}}{r_{i j}^{r_{i j}}}\right]
$$

## 20 years of CPU Optimization

- Single precision when possible
- Handtuned $1 /$ sqrt(x) instructions
- Single-instruction Multiple-Data
- Don't calculate zeros
- No charge on an atom? Don't do coulomh
- Interaction specific kernels
- Complex neighbor lists
- Fancy algorithms to extend timestep
- Bond constraints, virtual sites, etc.


## Stuff that might be hard to port to Cuda (I):

## Constraints

- $\Delta t$ limited by fast motions - 1fs
- Remove bond vibrations
- SHAKE (iterative, slow) - 2fs
- Problematic in parallel (won't work)
- Compromise: constrain h-bonds only 1.4fs
- GROMACS (LINCS):
- LINear Constraint Solver
- Approximate matrix inversion expansion
- Fast \& stable - much better than SHAKE
- Non-iterative
- Enables 2-3 fs timesteps
- Parallelizes (in theory at least)


## LINCS:,$t=2$ <br>  <br> A) Move w/o constraint


B) Project out motion along bonds
C) Correct for rotational extension of bond

## Stuff that might be hard to port to Cuda (2):

## Virtual sites

- Next fastest motions is H -angle and rotations of $\mathrm{CH}_{3} / \mathrm{NH}_{2}$ groups
- Try to remove them:
- Ideal H position from heavy atoms.
- $\mathrm{CH}_{3} / \mathrm{NH}_{2}$ groups are made rigid
- Calculate forces, then project back onto heavy atoms
- Integrate only heavy atom positions, reconstruct H's
- Enables 5fs timesteps!



## CPU Parallelization



Domain decomposition


Complex Load balancing

Lots of book-keeping

CPU limit today: 50-150 ns/day

Hard to decrease \# atoms / CPU


## Molecular Dynamics

 on Stream Processors
## Stream computing



## Our first cards were less fancy!

First Gromacs GPU project in 2002 with Ian Buck \& Pat Hanrahan, Stanford Promise of theoretical high FP performance on GeForces Severe limitations in practice...

## essons:

It's easy to achieve speedup relative to a slow reference implementation

Much harder to beat well-optimized CPU code that can use >1 core Don't write CPU code for GPUs

Find new optimizations instead

## Molecular Dynamics on GPUs



Focus on the code where we spend all the cycles on the CPU

## Bottlenecks



## OpenMM

Not sufficient to accelerate nonbonded interactions - need to send to/from GPU

Need to do entire simulation on GPU?
Not fun to rewrite 2M lines-of-code
 in a separate CUDA-Gromacs...

OpeniMiM: Core MD functionality in separate library Stanford (Pande), Stockholm (Us), Nvidia \& AMD

Fully public API, hardware-agnostic, use anywhere

Peter Eastman, Scott Legrand

## Nonbonded Interactions




- Divide into short-range / long-range
- Calculate short-range analytically
- Use approximations for the long-range


# O(N2) Algorithm on GPUs 

- An efficient algorithm should: - minimize global memory access
- avoid thread synchronization
- take advantage of symmetry


## O(N2) Algorithm on GPUs

- Group atoms into blocks of 32
- Interactions divide into $32 \times 32$ tiles
- Each tile is processed by a group of 32 threads
- Load atom coordinates and parameters into shared memory
- Each thread computes interactions of one atom with 32
 atoms
- Use symmetry to skip half the tiles


## $O\left(N^{2}\right)$ Algorithm on GPUs

- Each thread loops over atoms in a different order
- Avoids conflicts between threads
- No explicit synchronization needed
-Threads in a warp are always synchronized



## O(N) Algorithm on GPUs

## Hard. Why?

- Traditional O(N) methods (e.g. neighbor lists) slow on GPUs
- Out of order memory access

```
for i = 1 to numNeighbors
    load coordinates and parameters for neighbor[i]
    compute force
    store force for neighbor[i]
```

- Inner loop contains non-coalesced memory access


## Approach 1: Voxels

- Divide space into smaller voxels
- Compute all-vs-all interactions between adjacent voxels
- LAMMPS, NAMD
- Very efficient when all particles are identical
- Book-keeping can get expensive for complex systems



## Approach 2: Tiles

- Start with the $\mathbf{O}\left(\mathbf{N}^{2}\right)$ algorithm
- Exclude tiles with no interactions
- Like a neighbor list between blocks of 32 atoms

Problem 1: How do you construct good tiles?

Problem 2: How do you keep them good?


## Finding tiles with interactions

- Compute an axis aligned bounding box for the 32 atoms in each block
- Calculate the distance between boxes



## Keeping track of water

-Solvent molecules must be ordered to be spatially coherent

- Or bounding boxes will be very large
- Arrange along a space filling curve
- Reorder every ~100 time steps

Swapping is easy: Same parameters!


## Performance?

- Much faster than O(N²) for large systems
- Performance scales linearly

But

- Computes many more interactions than really required
- Computes all 1024 interactions in a tile, even if few/none are within the cutoff


## Finer-grained neighborlist

- For each tile with interactions:
- Compute distance of each atom in one block from the bounding box of the other block
- Set a flag for each atom



## Tile Force Computation

for $\mathrm{i}=1$ to 32
if (hasInteractions[i])
compute interaction with atom i

- All 32 threads must loop over atoms in the same order
-Requires a reduction to sum the forces
-For a few atoms, this is still much faster
-For many atoms, better to just compute all interactions


## PME (Long range coulomb)



1. Calculate scaled fractionals
2. Calculate B-spline coefficient
3. "smear" the charges over the grid points from spline-coefficient
4. Execute forward FFT
5. Calculate the reciprocal energy
6. Execute backward FFT
7. Calculate force gradient

## PME (Long range coulomb)

## CPU <br> spreading of charges



GPU
gathering of charges


- sort the atoms before gather !


## Gromacs \& OpenMM in practice

- GPUs supported in Gromacs 4.5 mdrun ... -device "OpenMM:Cuda"
- Same input files, same output files: "It just works"
- Subset of features work on GPUs for now (checked)
- No shortcuts taken on the GPU:
- At least same accuracy as on the CPU (<1e-6)
- Potential energies calculated, free energy works
- Prerelease availability: NOW! www.gromacs.org/gpu

Fermi (C20) performance over C10
BPTI ( $\sim 21 \mathrm{l}$ atoms)
Villin (600 atoms, implicit)

PME
Implicit

Reaction-field All-vs-all
$\qquad$

GPU performance over x86 CPU


NB: 2fs steps. GPUs still can't do vsites \& 5fs like CPUs

## Limitations

- Still hard to use long time steps on GPUs
- Virtual sites don't work
- We don't think you should go higher than 2fs steps without them
- Many of the CPU "tricks" can easily be written in Cuda, but we would end of with lots of kernels that must be called iteratively
- Hard to get multi-node GPU code to beat CPUs
- In the high end, we're all bandwidth-limited


## Hardware Caveats

- Beware of Memory Errors: happens on all hardware-...
- Gromacs runs tests to check for GPU memory errors
- Low-end consumer cards can sometimes be bad
- Even fine cards can exhibit random errors
- For production scientific work you might want Tesla-class Fermi cards...
- Why? ECC memory! (C2050/C2070)


## GPU weak scaling

Water box on a Nvidia Telsa2050


## More GPU/CPU comparisons

Nvidia Tesla 2050 vs Intel Core i7 920 (2.66 GHZ, 4 threads)

real space cost $\sim r_{c}^{3}$
FFT cost $\sim$ spacing $^{-3}$
for constant accuracy: spacing $=r_{c} / a$
total cost: $C_{\mathrm{pp}} r_{c}^{3}+C_{\mathrm{FFT}} a^{3} r_{c}^{-3}$

## Tiling circles is difficult

serial computing

stream computing


1. You need lots of cubes to cover a sphere
2. All interactions beyond cut-off must be zero

## The Art of Calculating Zeros



## Parallelize

the Problem


## These will soon be small computers

2022: ~300M cores
2020:~100M cores
2018:~30M cores
2016:~10M cores
2014:~3M cores
2012: ~1 M cores 2010:~300,000 cores

# We're all doing Embarrassing Parallelism But not the way you think. 

We're investing huge efforts in parallelizing algorithms that only reach 50-75\% scaling efficiency on large problems

Not a chance they will scale to 1B cores
Close-to-useless for smaller problems of commercial interest
100\% focus on programs, forget the problem we're solving
Ask taxpayers to foot the bill

# Scaling as an Obsession? 



## 

Gromacs has scaled to 150k cores on Jaguar @ ORNL

Only gigantic systems scale - limited number of applications


1M-100M atoms
But: Small systems won't scale to large numbers of cores!

## Ensemble Simulation



## Membrane Protein

 Insertion Free EnergyEvery dot is a simulation!


## Vesicle fusion - 1.5 M atoms



## Long fusion trajectories

- Run on Infiniband cluster, ~250-500 cores
- 7 vesicle pairs fusing in 100-250ns
- 75-100\% POPE lipids
- 2 non-fusing vesicle pairs, 100 ns \& 500ns - 50\% POPE lipids
- Interesting circumstantial observations, but hard to draw conclusions from


## Committer Analysis

- Pick 20 conformations along fusion path
- Restart with 20 different random seeds
- 'Shooting trajectories'
- 8 microseconds of additional simulation
- Run on capacity cluster, 16-32 cores each
- Folding@Home as cluster scheduler
- Calculate fusion probabilities


## Stalk commitment



## Statistics from 1000's of runs



## Scale the Problem, not Runs



- Stream Computing is the future for all HPC
- We're doing statistical mechanics!
- No algorithm will parallelize 5000 degrees of freedom over 1 billion processors
- Parallelize in the problem domain instead
- Node efficiency becomes the key measure


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